

## Solutions 11

## Exercise 11.1.

1. We can apply Eisenstein criterion with  $p(x) = x$ . We check that  $p(x)|x$ ,  $p(x)|x^3$  but  $p^2(x) \nmid x$  and  $p \nmid 1$ .
2. The polynomial  $\alpha^3 + \alpha + 1$  does not possess any root over  $\mathbb{F}_2$ . Since it has degree 3, it is irreducible as any decomposition would contain a degree 1 factor. Thus  $\alpha$  allows us to define a degree 3 extension, i.e the field  $\mathbb{F}_8 = \mathbb{F}_2[\alpha]$ . Besides let  $\omega$  be such that  $\omega^2 + \omega + 1 = 0$ . We have as usual  $\mathbb{F}_4 = \mathbb{F}_2[\omega]$ . Caveat :  $\mathbb{F}_4$  is never a subfield of  $\mathbb{F}_8$ . In general  $\mathbb{F}_{p^k} \subset \mathbb{F}_{p^{k'}}$  if and only if  $k$  divides  $k'$ . Now we explore case by case the number of points by fixing the value of  $x$ . Note that it is not necessary to try to solve all equation involved. You can use the Frobenius to reduce the number of computation. You can also observe that if  $(a, b)$  is a point on the curve, then  $(b^{-1}, a/b)$  and  $(b/a, a^{-1})$  are also points on the curve (if defined).

$x$	$\mathbb{F}_2$	$\mathbb{F}_4$	$\mathbb{F}_8$
0	0	0	0
1	$\emptyset$	$\emptyset$	$\alpha, \alpha^2, \alpha^4$
$\omega$		$\omega^2$	
$\alpha$			$1, \alpha, \alpha^6$
$\alpha^{-1}$			$\alpha^3, \alpha^4, \alpha^6$

So we get  $|\mathcal{K}_4(\mathbb{F}_2)| = 1$ ,  $\mathcal{K}_4(\mathbb{F}_4) = 1 + 2 \cdot 1 = 3$  and  $|\mathcal{K}_4(\mathbb{F}_8)| = 1 + 3 + 3 \cdot 3 + 3 \cdot 3 = 22$ .

3. To apply the theorems seen during the lecture, we need to change a bit the defining polynomial  $f$  in order to have a polynomial of the form  $x^4 + f_1(x, y)$  where the partial degree of  $f_1$  is  $\leq 3$ . This can be achieved by replacing  $y$  by  $x + y$ , since  $f(x, x + y) = x^4 + (y + 1)x^3 + x^2y + (y^2 + 1)x + y^3$ . This only changes the expression of the point of curve but not the parametres that we obtain. We consider the codes  $C(\mathcal{K}_4(\mathbb{F}_8), m)$  for  $4 \leq m \leq 6$  that yield  $[22, 4m - 6, \geq 26 - 4m]_8$ -codes as wanted. The table on the Internet shows that best known  $[22, 10]_8$ -code has minimum distance 10, best known  $[22, 14]_8$ -code distance 7 and best known  $[22, 14]_8$ -code distance 2.

## Exercise 11.2.

1. We notice that  $x^9 + 1$  has 1 as simple root and that  $x + 1$  is an irreducible polynomial that does not divide 1. So by Eisenstein criterion,  $f$  is irreducible.
2. One can simply check that 3 divides 6 or recall that  $\mathbb{F}_8$  is the set of roots of  $x^8 - x$  in any extension of  $\mathbb{F}_2$ . Now,  $\mathbb{F}_{64}^\times$  is a group of order 63 so it contains all 7-th roots of 1. We note that  $x \mapsto x^9$  is a multiplicative group homomorphism on  $\mathbb{F}_8^\times$ . Its kernel is the set of 9th roots of unity. Since  $\mathbb{F}_8^\times$  has order 7,  $\mathbb{F}_8$  does not contain any such root except 1, so the map is an isomorphism of  $\mathbb{F}_8^\times$ . It is clear that it remains a bijection on  $\mathbb{F}_8$ . On the other hand, on  $\mathbb{F}_{64}^\times$ ,  $x \mapsto x^9$  has a kernel of cardinality 9, namely  $\{\alpha^{7i}, 0 \leq i \leq 8\}$  where  $\alpha$  is a primitive element of  $\mathbb{F}_{64}$ . Besides, its image is in the set of the 7th root of unity,

ie  $\mathbb{F}_8^\times \subseteq \mathbb{F}_{64}^\times$ . Now, for cardinality reason, the morphism must be an epimorphism. So we have a  $9 - 1$  map onto  $\mathbb{F}_8^\times$ .

3. On  $\mathbb{F}_2$ , there are two points  $(0, 1)$  and  $(1, 0)$ . On  $\mathbb{F}_8$ , for any choice of  $x$  there is exactly a choice of  $y$ , because of the bijection property. So  $|\mathcal{F}_9(\mathbb{F}_8)| = 8$ . Now on  $\mathbb{F}_{64}$ , either  $x$  is one of the nine 9th roots of unity and  $y$  is 0 or  $x$  is one of the 55 non-9th root of unity, and  $y$  can take 9 values. In total we have  $|\mathcal{F}_9(\mathbb{F}_{64})| = 9 + 55 \cdot 9 = 504$  values. Weil bounds shows that any way  $-383 \leq |\mathcal{F}_9(\mathbb{F}_{64})| \leq 513$ . Now actually Weil bound applies projective curves, this curve is maximal if you consider the projective curve associated with it : consider the homogenised equation  $z^9 f(x/z, y/z) = x^9 + y^9 + z^9$  and count the number of non zero solutions up to homothety. You will find the 504 affine points that we have already (with  $z = 1$ ) and 9 additional points (with  $z = 0$ ). So there are 513 projective points on the curve which matches with Weil bound.

### Exercise 11.3.

1. We want to find the number of points  $(x, y, z) \in \mathbb{F}_{q^2}^3$  such that

$$\begin{aligned} z^{q+1} &= y^q + y \\ y^{q+1} &= x^q + x. \end{aligned}$$

Recall that the map  $x \mapsto x^q + x$  is the *trace* of  $\mathbb{F}_{q^2}$  over  $\mathbb{F}_q$  and is a surjective homomorphism, and that the map  $x \mapsto x^{q+1}$  is the *norm*, which means that it maps  $\mathbb{F}_{q^2}$  to  $\mathbb{F}_q$ . Now  $z$  can take  $q^2$  possible values, for each of which  $z^{q+1}$  is in  $\mathbb{F}_q$ . The trace map being a surjective homomorphism, we know that it maps to each value in  $\mathbb{F}_q$   $q$  times, so that there are  $q$  corresponding values of  $y$  such that  $y^q + y = z^{q+1}$ . But for each such value of  $y$ , there are again  $q$  values of  $x$  such that  $x^q + x$  maps to  $y^{q+1} \in \mathbb{F}_q$ . In total, there are  $q^4$  common zeros of  $f$  and  $g$ .

2. Consider  $f(x, y) = x^q + x - y^{q+1}$ . The space of polynomials of degree  $< m$  in  $\mathbb{F}_{q^2}[x, y]/(f)$  can be represented by

$$\Gamma_{< m} = \mathbb{F}_q[x]_{< m} \oplus y\mathbb{F}_q[x]_{< m-1} \oplus \cdots \oplus y^q\mathbb{F}_q[x]_{< m-q}$$

and is of dimension

$$\dim \Gamma_{< m} = (m-1)(q+1) - \frac{q(q-1)}{2} + 1 = 5m - 10$$

for  $q := 4$ . Now consider  $g(x, y, z) = y^q + y - z^{q+1}$ . The space of polynomials of degree  $< n$ , in  $\mathbb{F}_{q^2}[x, y, z]/(f, g)$  can be represented by

$$\Gamma_{< n} \oplus z\Gamma_{< n-1} \oplus \cdots \oplus z^q\Gamma_{< n-q},$$

for  $n > q$ . This space has dimension

$$\sum_{m=n-q}^n (5m - 10) = 25n - 100.$$

3. We consider the curve given by  $f(x, y, z) = 0, g(x, y, z) = 0$  and define the corresponding AG code  $C = \text{Im}\phi$ , where

$$\begin{aligned} \phi : \Gamma_{<n} \oplus z\Gamma_{<n-1} \oplus \cdots \oplus z^q\Gamma_{<n-q} &\rightarrow \mathbb{F}_q^{q^4} \\ h &\mapsto (h(a, b, c) : f(a, b, c) = g(a, b, c) = 0). \end{aligned}$$

By Bézout's theorem, we know that the number of common zeros of any  $h$  and the curve is upper bounded by  $\deg(h)\deg(f)\deg(g) < n(q+1)^2 = 25n$  (note that  $h$  belongs to such a space that it cannot be expressed as  $h = h_1f + h_2g$  for some polynomials  $h_1$  and  $h_2$ , so that Bézout's theorem can be applied).

We would like to make sure that  $\phi$  is injective. For that, it is enough to impose the condition  $n(q+1)^2 < q^4$ , i.e.,  $n \leq 10$ , which ensures that for any  $h$  in the domain,  $\phi(h)$  is not all zeros. We then get the following parameters for the code:

$$\begin{aligned} \dim C &= 25n - 100 \\ \text{dist} C &\geq q^4 - n(q+1)^2 = 256 - 25n. \end{aligned}$$

#### Exercise 11.4.

1. Suppose that  $n = 1$ , so that  $P(x_1)$  is a univariate polynomial of degree  $d$  over  $\mathbb{F}_q$ . Then we know that  $P$  has at most  $d$  roots in  $\mathbb{F}_q$ , therefore for  $x_1$  chosen uniformly at random in  $\mathbb{F}_q$ , we get

$$\Pr[P(x_1) = 0] \leq d/q.$$

2. Now suppose that for a given  $n \geq 2$ , for any nonzero  $(n-1)$ -variate polynomial  $P(x_1, \dots, x_{n-1})$  of degree  $d$ ,  $\Pr[P(x_1, \dots, x_{n-1}) = 0] \leq d/q$  when  $x_1, \dots, x_{n-1}$  are chosen uniformly in  $\mathbb{F}_q$ . We want to prove that for a nonzero  $n$ -variate polynomial  $P(x_1, \dots, x_n)$  of degree  $d$ ,  $\Pr[P(x_1, \dots, x_n) = 0] \leq d/q$  when  $x_1, \dots, x_n$  are chosen uniformly in  $\mathbb{F}_q$ .

Write

$$\begin{aligned} P(x_1, \dots, x_n) &= \sum_{i=0}^d x_1^i P_i(x_2, \dots, x_n) \\ &= \sum_{i=0}^j x_1^i P_i(x_2, \dots, x_n), \end{aligned}$$

where  $j$  is the largest index such that  $P_j(x_2, \dots, x_n)$  is not the zero polynomial. Note that  $\deg P_j + j \leq d$ . We have that

$$\begin{aligned} \Pr_{x_1, \dots, x_n} [P(x_1, \dots, x_n) = 0] &= \Pr_{x_1, \dots, x_n} [P(x_1, \dots, x_n) = 0 \wedge P_j(x_2, \dots, x_n) = 0] \\ &\quad + \Pr_{x_1, \dots, x_n} [P(x_1, \dots, x_n) = 0 \wedge P_j(x_2, \dots, x_n) \neq 0]. \end{aligned}$$

But

$$\Pr_{x_1, \dots, x_n} [P(x_1, \dots, x_n) = 0 \wedge P_j(x_2, \dots, x_n) = 0] \leq \Pr_{x_2, \dots, x_n} [P_j(x_2, \dots, x_n) = 0] \leq \frac{\deg P_j}{q},$$

where we have used the induction hypothesis. Also,

$$\Pr_{x_1, \dots, x_n} [P(x_1, \dots, x_n) = 0 \wedge P_j(x_2, \dots, x_n) \neq 0] \leq \Pr_{x_1} [P(x_1, \dots, x_n) = 0 | x_2, \dots, x_n \text{ are s.t. } P_j(x_2, \dots, x_n) \neq 0].$$

But when we condition on the event that  $P_j(x_2, \dots, x_n) \neq 0$ , we have that  $P(x_1, \dots, x_n) = \sum_{i=0}^j x_1^i P_i(x_2, \dots, x_n)$  is a **nonzero** degree- $j$  univariate polynomial in  $x_1$ . We can now apply the result from part 1 and deduce that the probability that such a polynomial evaluates to 0 is upper-bounded by  $j/q$ .

We thus have that

$$\Pr[P(x_1, \dots, x_n) = 0] \leq \frac{\deg P_j + j}{q} \leq d/q,$$

which completes the induction step.

3. By definition,

$$\Pr[P(x_1, \dots, x_n) = 0] = \frac{\# \text{ roots } (a_1, \dots, a_n) \text{ of } P}{q^n}.$$

Putting this together with the result of part 2, we conclude that the number of roots  $(a_1, \dots, a_n)$  of  $P$  over  $\mathbb{F}_q$  is less than or equal to  $dq^{n-1}$ .

Note that we could consider the alternative problem of bounding the number of roots  $(a_1, \dots, a_n)$  of  $P$  where all  $a_i$  belong to a subset  $S \subseteq \mathbb{F}_q$ . Then we could follow a similar method to upper bound the number of such roots by  $d|S|^{n-1}$ .

### Exercise 11.5.

1. Write  $f(x, y)$  as  $\sum_i f_i(y)x^i$ . For any  $\beta \in I$ ,  $f(x, \beta) = \sum_i f_i(\beta)x^i$  is a univariate polynomial in  $x$  of degree  $< k$ , and it has at least  $k$  roots ( $x$  instantiated with all elements of  $I$ ). Hence it is identically zero, and  $f_i(\beta) = 0 \forall i$ . Now for each such index  $i$ , we thus have that  $f_i(\beta) = 0$  for all elements  $\beta$  of  $I$ , i.e.,  $f_i(y)$  has at least  $k$  roots. But it is a univariate polynomial in  $y$  of degree  $< y$ , so that  $f_i(y)$  must be identically zero. We thus have that  $f(x, y) = \sum_i f_i(y)x^i$  is identically 0.
2. The code  $C$  is the image of the map

$$\begin{aligned} \phi : \mathbb{F}_q[x, y]_{<k, <k} &\rightarrow \mathbb{F}_q^2 \\ f &\mapsto (f(a, b) : (a, b) \in \mathbb{F}_q^2) \end{aligned}$$

We would like to show that  $\phi$  is injective. Take an element  $f$  of  $\ker \phi$ .  $f$  is such that it evaluates to 0 on all  $(a, b) \in \mathbb{F}_q^2$ . Apply part 1 with  $I = \mathbb{F}_q$  to get that  $f$  is the zero polynomial. Thus  $\phi$  is injective and  $\dim C = \dim(\text{Im} \phi) = k^2$ .

3. Finding such a nonzero polynomial  $Q(x, y, z) = \sum_{i=0}^{\ell-1} \sum_{j=0}^{\ell-1} \sum_{k=0}^{t-1} q_{ijk} x^i y^j z^k$  amounts to finding its  $\ell^2 t$  coefficients. The conditions  $Q(\alpha, \beta, \gamma_{\alpha, \beta}) = 0$  for all  $\alpha, \beta \in \mathbb{F}_q$  correspond to  $q^2$  linear equations that the  $\{q_{ijk}\}$  must satisfy. As long as  $\ell^2 t > q^2$ , this linear system has a nontrivial solution.