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COMPRESSION FOR THE AWGN CHANNEL

$$G = \begin{bmatrix} \\ \\ \end{bmatrix}_{k \times n} \quad n < k$$

Choose each column of G at random from set of all binary vectors of weight d . What is the prob that $\text{rank}(G) = n$ as $n, k \rightarrow \infty$?

SASON - WIECHMAN

$$H = \begin{bmatrix} \\ \\ \end{bmatrix}_{(n-k) \times n}$$

Need the fraction of rows of weight k , denote. Consider the Kolchin ensemble we used for the d ones are distributed uniformly at random among $(n-k)$ bins in each column. What is the prob that an number of ones landed in the first bin?

$$P(\text{odd \# in 1st bin}) = \binom{1}{n-k} \left(\frac{n-k-1}{n-k}\right)^{d-1} + \binom{1}{n-k}^3 \left(\frac{n-k-1}{n-k}\right)^{d-3} + \dots + \binom{1}{n-k}^d \quad \{d \text{ odd}\}$$

$$\text{(OR)} \quad \left(\frac{1}{n-k}\right)^{d-1} \left(\frac{n-k-1}{n-k}\right) \quad \}$$

Geometric progression $\frac{a(1-r^n)}{1-r}$

$$a = \binom{1}{n-k} \left(\frac{n-k-1}{n-k}\right)^{d-1} \quad r = \left(\frac{1}{n-k}\right)^2 \left(\frac{n-k-1}{n-k}\right)$$

$$n = \left\lceil \frac{d}{2} \right\rceil \quad = \frac{1}{(n-k-1)^2}$$

29986359
 42031088
 9840766538

2 → 2
 3 → 4
 4 → 4

$P(\text{odd \# in 1st bin})$

$$= \binom{1}{n-k} \binom{n-k-1}{n-k}^{d-1} \frac{\left(1 - \frac{1}{(n-k-1)^{2\lceil d/2 \rceil}}\right)}{(n-k-1)^2 - 1} \cdot (n-k-1)^2$$

$$= \frac{(n-k-1)^{d+1 - 2\lceil d/2 \rceil}}{(n-k)^d} \cdot \frac{(n-k-1)^{2\lceil d/2 \rceil} - 1}{(n-k-1)^2 - 1}$$

$\triangleq P_1$

$P(\text{row of weight } j) = \binom{n}{j} P_1^j (1 - P_1)^{n-j} = \Gamma_j$

$P_1 \sim O\left(\frac{1}{n}\right)$ $\left\{ \begin{array}{l} \text{Intuitively, as } n \rightarrow \infty, \text{ most entries will} \\ \text{have zero 1s, which is an even number.} \\ \text{Hence we should expect } P_1 \rightarrow 0 \text{ as } n \rightarrow \infty \end{array} \right.$

$\Gamma(x) = \sum_j \Gamma_j x^j$

$= \sum_{j=0}^n \binom{n}{j} P_1^j (1 - P_1)^{n-j} x^j$

$= [1 - P_1 + P_1 x]^n$ $\left\{ \begin{array}{l} \text{check: } \Gamma(1) = \sum_j \Gamma_j = 1 \checkmark \end{array} \right.$

$\frac{H(x|Y)}{n} \geq 1 - C - (1-R) \left(1 - \frac{1}{2 \ln 2} \sum_{p=1}^{\infty} \frac{\Gamma(p)}{p(2^p - 1)}\right)$

As $n \rightarrow \infty$, we have that $\Gamma(x) > 0$ iff

$1 - P_1 + P_1 x \geq 1$

$\Leftrightarrow x \geq 1$

$g_p \triangleq \int_0^{\infty} \underbrace{a(\ell)}_{\text{cond pdf of LLR given input zero}} (1 + e^{-\ell}) \tanh^{2p}\left(\frac{\ell}{2}\right) d\ell$

$$\int_{-\infty}^{\infty} \frac{1}{\sqrt{2\pi}} e^{-\frac{x^2}{2}} dx = 1$$

BIANGAN :

$$a(l) = \sqrt{\frac{\sigma^2}{8\pi}} e^{-\left[\frac{(l - \frac{\mu}{\sigma^2})^2 \sigma^2}{8}\right]}$$

From Richardson,
Shokrollahi & Ural
Example 8

$$g_p = \int_0^{\infty} a(l) \tanh^{2p}\left(\frac{l}{2}\right) dl + \int_0^{\infty} a(l) e^{-l} \tanh^{2p}\left(\frac{l}{2}\right) dl$$

$l' = -l$

$$\int_0^{\infty} a(-l') e^{l'} \tanh^{2p}\left(\frac{-l'}{2}\right) dl'$$

$$= \int_{-\infty}^0 a(l) \tanh^{2p}\left(\frac{l}{2}\right) dl$$

(from symmetry p)

$$g_p = \int_{-\infty}^{\infty} a(l) \tanh^{2p}\left(\frac{l}{2}\right) dl \quad (\text{for symmetric distri})$$

$$= \int_{-\infty}^{\infty} \sqrt{\frac{\sigma^2}{8\pi}} e^{-\left[\frac{(l - \frac{\mu}{\sigma^2})^2 \sigma^2}{8}\right]} \left[\frac{e^l - 1}{e^l + 1}\right]^{2p} dl \quad \left\{ \begin{array}{l} \because \tanh x \\ = \frac{e^{2x} - 1}{e^{2x} + 1} \end{array} \right.$$

$$g_p = \int_0^{\infty} a(l) (1 + e^{-l}) \left[\tanh\left(\frac{l}{2}\right)\right]^{2p} dl$$

$$\leq \int_0^{\infty} a(l) (1 + e^{-l}) dl = 1$$

$\{\because \tanh x \leq 1\}$

Hence we have that

$$\frac{H(X|Y)}{n} \geq X - C = (X - R) = R - C$$

$$I(X; Y) = H(X) - H(X|Y)$$

$$\leq nR - n(R - C)$$

$$= nC$$

As far as this bound is concerned, the LDPC encoder seems optimal.