

$$W = \frac{2^{k-2}}{N} G^T \left\{ I_k + \begin{bmatrix} 1 & \dots & 1 \\ \vdots & & \vdots \\ 1 & \dots & 1 \end{bmatrix}_{k \times k} \right\} G_{k \times n}$$

$$= \frac{2^{k-2}}{N} G^T G + \frac{2^{k-2}}{N} \underbrace{d}_{\substack{G^T \\ n \times k}} \begin{bmatrix} 1 & \dots & 1 \\ \vdots & & \vdots \\ 1 & \dots & 1 \end{bmatrix}_{k \times n}$$

$$\underbrace{\begin{bmatrix} 1 & \dots & 1 \\ \vdots & & \vdots \\ 1 & \dots & 1 \end{bmatrix}}_{n \times n}$$

$$W_{ij} = \frac{2^{k-2}}{N} [g_i^T g_j + d]$$

$$W = \frac{1}{N} \left[ 4 \Xi_2^T \Xi_2 - 2 \Xi_2^T \underline{1} - 2 \underline{1}^T \Xi_2 + \underline{1}^T \underline{1} \right]$$

Need to take care of real vs.  $GF_2$  operations.

Codeword generation:

Define  $\Xi = A G$   
↑  
real ops.

Actual codewords =  $\Xi \pmod{2}$

Actual patterns =  $2 \Xi \pmod{2} - \underline{1} \equiv 2 \Xi_2 - \underline{1}$

$$\Xi_2^T \Xi_2 = (G^T A^T)_{\pmod{2}} (A G)_{\pmod{2}}$$

$$\neq \underbrace{(G^T A^T A G)}_{\text{all real ops.}} \pmod{2}$$

$$= (G^T G) \pmod{2}$$

~~$(A^T A)_{\pmod{2}} = \underline{0} \Rightarrow \Xi_2^T \Xi_2 = \underline{0}$~~

$$\begin{matrix} \underline{1}^T & \underline{1} \\ \hline n \times 2^k & 2^k \times n \end{matrix} = 2^k \underline{1}_{n \times n}$$

$$\begin{matrix} \underline{1}^T & \underline{1} \\ \hline n \times 2^k & 2^k \times n \end{matrix} = 2^{k-1} \underline{1}_{n \times n} = \underline{1}^T \Xi_2$$

$$W = \frac{1}{N} \left[ 4 \Xi_2^T \Xi_2 - 2^{k+1} \mathbb{1}_{N \times N} + 2^k \mathbb{1}_{N \times N} \right]$$

$$= \frac{1}{N} \left[ 4 \Xi_2^T \Xi_2 - 2^k \mathbb{1}_{N \times N} \right]$$

Suppose that  $W = \Gamma^T \Delta \Gamma$

↓  
should map vectors of small Euclid norm to vectors of small  $\infty$ -norm

Check to see: Sum a few cols of the Hadamard matrix, see what the obtained  $\infty$ -norm is

### MAC DISTORTION EXPONENTS:

Informed tx upper bound, a closer look:

Equivalent MIMO system

$$b' = \frac{T}{|K|S} = \frac{b}{|K|}$$

Rate of transmission  $R' = |K| R(H)$

Distortion obtained for the MIMO equivalent

$$D'(P) = \mathbb{E} \left[ e^{-b' R'} \right] \quad \left\{ \begin{array}{l} \text{Source-channel separation is} \\ \text{optimal here} \end{array} \right.$$

$$= \mathbb{E} \left[ e^{-b R(H)} \right]$$

Distortion for the actual system is at least equal to of the MIMO system:

$$D(P) \geq D'(P) = \mathbb{E} \left[ e^{-b R(H)} \right]$$

$$\geq \mathbb{E} \left[ e^{-b \cdot \frac{1}{|K|} \log \det(\cdot)} \right] \quad \forall |K|$$

## MODIFYING THE HOPFIELD METRIC

$$W_{ij} = \frac{1}{N} \sum_{m=1}^M \xi_i^m \xi_j^m = \frac{1}{N} \Xi^T \Xi \triangleq W$$

$$\downarrow$$

$$\frac{1}{N} \sum_{m=1}^M \Omega_m \xi_i^m \xi_j^m = \frac{1}{N} \Xi^T \begin{bmatrix} \Omega_1 & & \\ & \Omega_2 & \\ & & \ddots \\ & & & \Omega_M \end{bmatrix} \Xi \triangleq W$$

$$= \begin{bmatrix} | & & | \\ \xi^1 & & \xi^M \\ | & & | \end{bmatrix} \begin{bmatrix} -\Omega_1 \xi^1 \\ -\Omega_2 \xi^2 \\ \vdots \\ -\Omega_M \xi^M \end{bmatrix}$$

What about  $\frac{1}{N} \Xi^T \Omega \Xi$  for  $\Omega$  not necessarily diagonal?

$$W_{ij} = \frac{1}{N} \sum_{k,l} (\Xi^T)_{ik} \Omega_{kl} \Xi_{lj}$$

Recovery equations:

$$x_i^{(t+1)} = \max \left\{ \sum_{j \neq i} W_{ij} x_j^{(t)}, 0 \right\} = \max \left\{ 0, W_{ij} x_j^{(t)} \right\}$$

Suppose that

Diagonal entries of  $W'$ :

$$W'_{ii} = \sum_{m=1}^M \Omega_m \left( \xi_i^m \right)^2 = \sum_{m=1}^M \Omega_m$$

$$W_N = W' - \left( \sum_{m=1}^M \Omega_m \right) \mathbf{I}_N$$

Suppose we feed in a pattern  $\xi'$ :

$$W_N \xi' = W' \xi' - \frac{1}{N} \left( \sum_{m=1}^M \Omega_m \right) \xi'$$

Assume that the  $\xi^m$  are (almost) orthogonal.

$$W_N \xi' = \frac{1}{N} \xi' (\Omega_1, \dots, \Omega_M) - \frac{1}{N} \left( \sum_{m=1}^M \Omega_m \right) \xi'$$

$$= \xi' \left[ -\Omega_1, -\frac{\sum_{m=1}^M \Omega_m}{N} \right]$$

{ perfectly orthogonal, as  $M=N$  in this case

$$E = \frac{k}{R}$$

$$n = \frac{k}{R}$$

$$2^{\frac{2k}{R}} < n$$

$$2^{\frac{2k}{R}} < n$$

$$Rn < \log_2 n$$

$$R < \frac{\log_2 n}{n}$$

$$\Omega_1 > \frac{\alpha M}{N}$$

$$N\Omega_1 > \alpha M$$

$$\Omega_1 > \frac{\alpha M}{N}$$

$$W_n \xi_1' = \frac{1}{N} \xi_1'(\Omega_1, N) + \frac{1}{N} \xi_2' \Omega_2 \langle \xi_2', \xi_1' \rangle + \frac{1}{N} \xi_3' \Omega_3 \langle \xi_3', \xi_1' \rangle$$

$$+ \dots + \frac{1}{N} \xi_M' \Omega_M \langle \xi_M', \xi_1' \rangle - \frac{\sum_{m=1}^M \Omega_m}{N} \xi_1'$$

$$= \xi_1' \left[ \Omega_1 - \frac{\sum_{m=1}^M \Omega_m}{N} \right] + \frac{1}{N} \sum_{m=2}^M \xi_m' \Omega_m \langle \xi_m', \xi_1' \rangle$$

would like this to be positive  
(N large helps, if all  $\Omega_i$  +ve)

would like keep the m of all cor small

Observations:

① Suppose that  $\sum_{m=1}^M \Omega_m \approx c$

Then  $\Omega_1 - \frac{\sum_{m=1}^M \Omega_m}{N} > 0$

$$\Leftrightarrow \Omega_1 > \frac{c}{N} \Rightarrow \sum_{m=1}^M \Omega_m > \frac{Mc}{N}$$

$$c > \frac{M}{N}$$

$$\boxed{N > M} \text{ if } c > 0$$

$$\boxed{M > N} \text{ if } c < 0$$

Gold sequences:

Length of the sequences  $n = 2^m - 1$   
 # of sequences =  $2^m + 1$  (known these are all cyclically d  
 max  $\{ | \text{corr. betn cyclic shifts of two Gold seqs.} | \} = 1 + \sqrt{2^{m-1}}$   
 $= 1 + \sqrt{2(n+1)}$

All these correlations are with sequences with norm 1  
 (this is trivially satisfied for binary seqs, being  $\xi_{\pm}$ )

# of seqs including cyclic shifts

$$= n(2^m + 1)$$

$$= n(n+2) = n^2 + 2n$$

What happens to the correlation term in this case?

$$\left| \frac{1}{N} \sum_{m=2}^M \xi_m \Omega_m \langle \xi_m, \xi_1 \rangle \right|_{M = N^2 + 2N} \quad \left\{ \begin{array}{l} 1:1 \text{ is done} \\ \text{componentwise} \\ (\infty\text{-norm?}) \end{array} \right.$$

$$\leq \frac{1}{N} \sum_{m=2}^{N^2+2N} \left| \xi_m \Omega_m \left[ 1 + \sqrt{2(N+1)} \right] \right|$$

$$\leq \frac{|\Omega_{\max}| \cdot (1 + \sqrt{2(N+1)})}{N} \cdot \sum_{m=2}^{N^2+2N} (1)$$

$$\approx \frac{|\Omega_{\max}|}{N} \cdot \sqrt{2N} \cdot N(N+1)$$

Case 1: Say  $\Omega_i = o(1) \quad \forall i$

$$\Omega_i \neq \Omega_j \quad \forall i \neq j$$

$$\Omega_i < 0 \quad \forall i$$

$$\Omega_1 = \frac{\sum_{m=1}^M \Omega_m}{N} \approx \frac{M o(1)}{N} = o(1)$$

$$= \frac{(N^2 + 2N) o(1)}{N} = o(1)$$

$$= (N+2) o(1) = o(1)$$