

Exercise Sheet 4

Exercise 4.1. Recall that, for integer parameters n, q , we defined Krawtchouk polynomials as

$$K_\ell(x) := \sum_{r=0}^{\ell} \binom{x}{r} \binom{n-x}{\ell-r} (-1)^r (q-1)^{\ell-r}.$$

First, show that

$$(1 + (q-1)z)^{n-x} (1-z)^x = \sum_{r=0}^x \binom{x}{r} (1 + (q-1)z)^{n-r} (-qz)^r.$$

Then, compute and compare the coefficients of z^ℓ on both sides. Using this, derive an alternative form of Krawtchouk polynomials.

Exercise 4.2. Let n, k, d be positive integers ($k \leq n - d + 1$), and consider the ensemble of all $(n-k) \times n$ matrices over \mathbb{F}_q of the form

$$H = (A \mid I),$$

where I is the $(n-k) \times (n-k)$ identity matrix and A is arbitrary, and define a probability distribution on this ensemble that is induced by assuming a uniform distribution over the $(n-k) \times k$ matrices A over \mathbb{F}_q .

1. Show that for every nonzero vector $y \in \mathbb{F}_q^n$,

$$\Pr[Hy^\top = 0] = \begin{cases} 0 & \text{if the first } k \text{ entries in } y \text{ are zero} \\ q^{k-n} & \text{otherwise.} \end{cases}$$

2. Show that

$$\Pr[H \text{ contains } d-1 \text{ dependent columns}] \leq \rho,$$

where

$$\rho := q^{k-n} \cdot \sum_{i=1}^{d-1} \left(\binom{n}{i} - \binom{n-k}{i} \right) (q-1)^{i-1}.$$

3. Deduce that all but a fraction at most ρ of the systematic linear $[n, k]$ codes over \mathbb{F}_q (i.e., codes with parity check matrices of the form above) have minimum distance at least d .

Exercise 4.3. Denote by $A(n, d, w)$ the maximal number of codewords in a binary code of length n and minimum distance at least d for which all codewords have the same weight w .

1. Let \mathcal{C} be an (n, k, d) -code containing the zero codeword. Show that the number of nonzero words in \mathcal{C} with weight up to d is at most $A(n, d, d)$.
2. Show that $A(n, 2k-1, w) = A(n, 2k, w)$.

3. Show that $A(n, 2k, w) \leq nA(n-1, 2k, w-1)/w$.
(*Hint:* Suppose that \mathcal{C} is a binary code of length n , minimum distance at least $2k$ for which all codewords have weight w . As in the proof of Plotkin bound, arrange the words of \mathcal{C} as rows of a matrix and upper bound the number of ones in each column of this matrix.)
4. Conclude that

$$A(n, 2k, w) \leq \left\lfloor \frac{n}{w} \left\lfloor \frac{n-1}{w-1} \left\lfloor \dots \left\lfloor \frac{n-w+k}{k} \right\rfloor \dots \right\rfloor \right\rfloor \right\rfloor.$$