

# Numerical experiments to design codes on Grassmannians

Who? Bertrand '

When? Algo Day – September 30th, 2009

# What is a Grassmanian ?

## Definition

Fix integers  $d \leq n$ .

The set of all  $d$ -dimensional vector subspaces of  $\mathbb{R}^n$  is called the **Grassmannian** and denoted  $\mathcal{G}_{n,d}$ .

## Properties

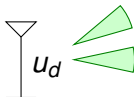
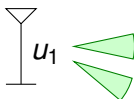
Any element can be represented by a  $n \times d$  matrix  $X$ , unique up to a left multiplication by  $\mathbf{GL}_d(\mathbb{R})$ .

- It is a *variety* ( $\rightsquigarrow$  we can do geometry).
- It is an *homogeneous space* ( $\rightsquigarrow$  under the action of  $\mathbf{GL}_n(\mathbb{R})$ , it looks everywhere the same).

# Where coding theory comes in

The setting

Multiple antennas.



Model :

$$Y = HX + W \quad \text{with } X = \begin{pmatrix} u_1 \\ \vdots \\ u_d \end{pmatrix}$$

Non coherent case :  $H$  is not known . Up to noise, the receiver can only “see” the space spanned by  $u_1, u_2 \dots u_d$ .

# A mathematical problem

**Definition** For  $x, x' \in \mathcal{G}_{n,d}$ , define the **product distance** by

$$d_{\Pi}(x, x') = \det(\mathbf{I}_n - p_x \circ p'_x)$$

with  $p_x = x(x^\dagger x)^{-1}x^\dagger$  the orthogonal projector on  $x$ .

**Problem** Given  $n, d$  and  $N$ , find  $C \subseteq \mathcal{G}_{n,d}$  such that

- $|C| = N$
- $\min_{x, x' \in C} d_{\Pi}(x, x')$  be as big as possible.

# A first attempt

**Goal** Find numerically good configuration, try to identify them, detect families.

cf. Work by Conway, Hardin & Sloane, *Packing Lines, Planes, etc., Packings in Grassmannian Spaces*, Experimental Mathematics, Vol. 5 (1996), 139-159, for other distances on  $\mathcal{G}_{n,d}$   
We focus on  $n = 4, d = 2$ .

## Method

- Pick randomly  $N$  points  $C$  of  $\mathcal{G}_{n,d}$ .
- Imagine that a “force” separate them if the distance is too small by trying to minimise a potential

$$\mathcal{P}(C) = \sum_{x, x' \in C} \frac{1}{(d_{\Pi}(x, x') - \alpha)^k}$$

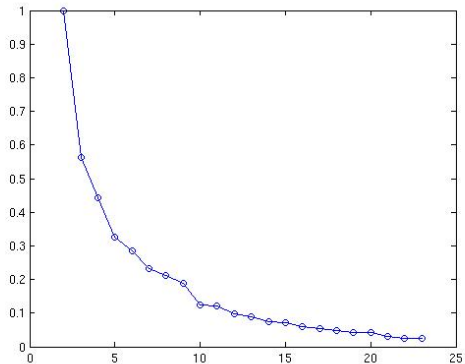
(for some parameter  $\alpha, k$  to adjust).

- Try to improve the minimum distance by small moves.

# The Hooke - Jeeves algorithm

- Start with a configuration  $C$
- [Exploration move ] Sample the potential function  $\mathcal{P}$  successively on each direction (in the tangent space) by a step  $\Delta$  to find a better configuration.
- [Pattern move] If the exploratory move gives a new configuration, use this direction to find a new configuration and repeat.
- Else, reduce the step  $\Delta$ .
- Stop when  $\Delta$  is small.

## Numerical result for $d_n$



## How to visualise the configuration (I)

- Compute the **principal angles**,  
i.e. find an orthonormal basis  $(u_1, u_2)$  of  $x \in C$  and an orthonormal basis  $(u'_1, u'_2)$  such that  $(u_1, u_2)$  has maximal angle in  $[0, \frac{\pi}{2}]$ .

By the way,  $d_{\Pi}(x, x') = \sin(u_1, u'_1)^2 \sin(u_2, u'_2)^2$

Are they nice values ? e.g.  $90^\circ$ ,  $45^\circ$ ,  $30^\circ$  ...



## How to visualise the configuration (II)

- Use **binocular code**,  
Represent a plane  $x = \langle u_1, u_2 \rangle$  of  $\mathcal{G}_{4,2}$  by the rotation  $\alpha = p_x - p_{x^\perp}$  in quaternions.

### Property

There exists  $\ell, r$  purely imaginary quaternions such that  $\alpha = x \mapsto \bar{\ell}xr$ . The pair  $\pm(\ell, r)$  is unique up to sign and equal to  $\pm(u_1\bar{u}_2 - u_2\bar{u}_1, \bar{u}_2u_1 - \bar{u}_1u_2)$ .

Thus we can visualise any code  $C$  as a subset of  $S^2 \times S^2$ .

Given  $(\ell, r)$  and  $(\ell', r')$  in  $S^2 \times S^2$ , the principal angles more or less are given by

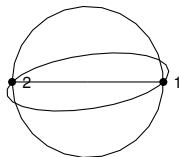
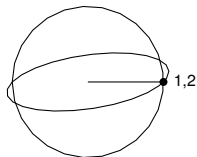
$$\frac{1}{2} \left( \widehat{(\ell, \ell')} \pm \widehat{(r, r')} \right).$$

## Results

■  $N = 2 : C = \left\{ \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \right\}$

Principal angles  $(\frac{\pi}{2}, \frac{\pi}{2})$

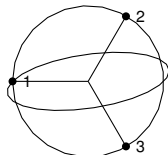
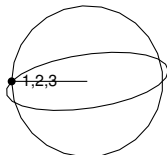
$$d_{\min} = 1$$



■  $N = 3 :$

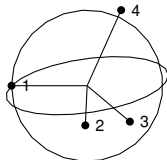
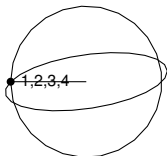
Principal angles :  $(\frac{\pi}{3}, \frac{\pi}{3})$

$$d_{\min} = \sin^4 \frac{\pi}{3} = \frac{9}{16}.$$



■  $N = 4$

Principal angles :  $(\theta, \theta)$  with  $\theta = \arcsin \frac{\sqrt{6}}{3}$ .

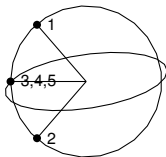
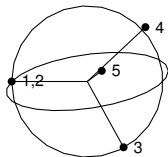


$d_{\min} = \frac{4}{9}$

■  $N = 5$

Principal angles :  $(49.1^\circ, 49.1^\circ)$  or  $(84.2^\circ, 35.1^\circ)$

$d_{\min} = \frac{241}{144} - \frac{7\sqrt{3}}{9} \simeq 0.3265$



- $N = 6$  Principal angles ( $47^\circ, 47^\circ$ ) or ( $90^\circ, 32^\circ$ )  
Minimal distance :  $\frac{17-4\sqrt{13}}{9} \simeq .2864$

